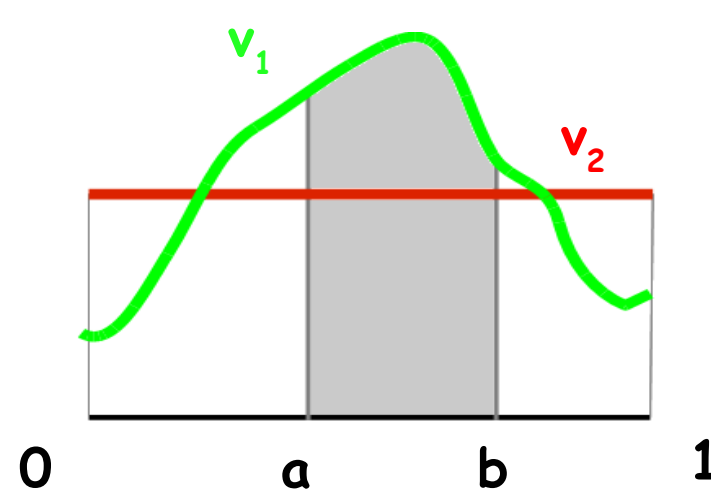


## Cake Cutting

Fundamental problem in fair division; models the allocation of a divisible resource (time, land, computer memory) among agents with heterogeneous preferences.

- The cake is the interval  $[0, 1]$
- Set of agents  $N = \{1, \dots, n\}$
- Each agent  $i$  has valuation function  $V_i$  over the cake, which is the integral of a value density function  $v_i$



- A piece of cake is a finite union of disjoint subintervals of  $[0,1]$ .
- The valuation of agent  $i$  for a piece  $X$  is given by the integral of their density function over the piece:

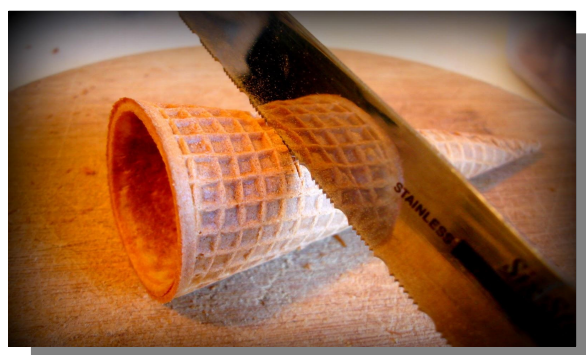
$$V_i(X) = \sum_{I \in X} \int_I v_i(x) dx$$

- An allocation  $A = (A_1, \dots, A_n)$  is an assignment of pieces to agents such that each agent  $i$  receives piece  $A_i$  and all the  $A_i$  are disjoint.
- Allocation  $A$  is **proportional** if  $V_i(A_i) \geq 1/n$ ,  $\forall i \in N$  and **envy-free** if  $V_i(A_i) \geq V_i(A_j)$ ,  $\forall i, j \in N$ .

## Query Model (Robertson & Webb)

- All the discrete cake cutting protocols interact with the players using two types of queries:
  - > **Cut** $(x, a)$ : Agent  $i$  returns  $y$  such that  $V_i([x, y]) = a$
  - > **Evaluate** $_i(x, y)$ : Agent  $i$  returns  $a$  such that  $V_i([x, y]) = a$

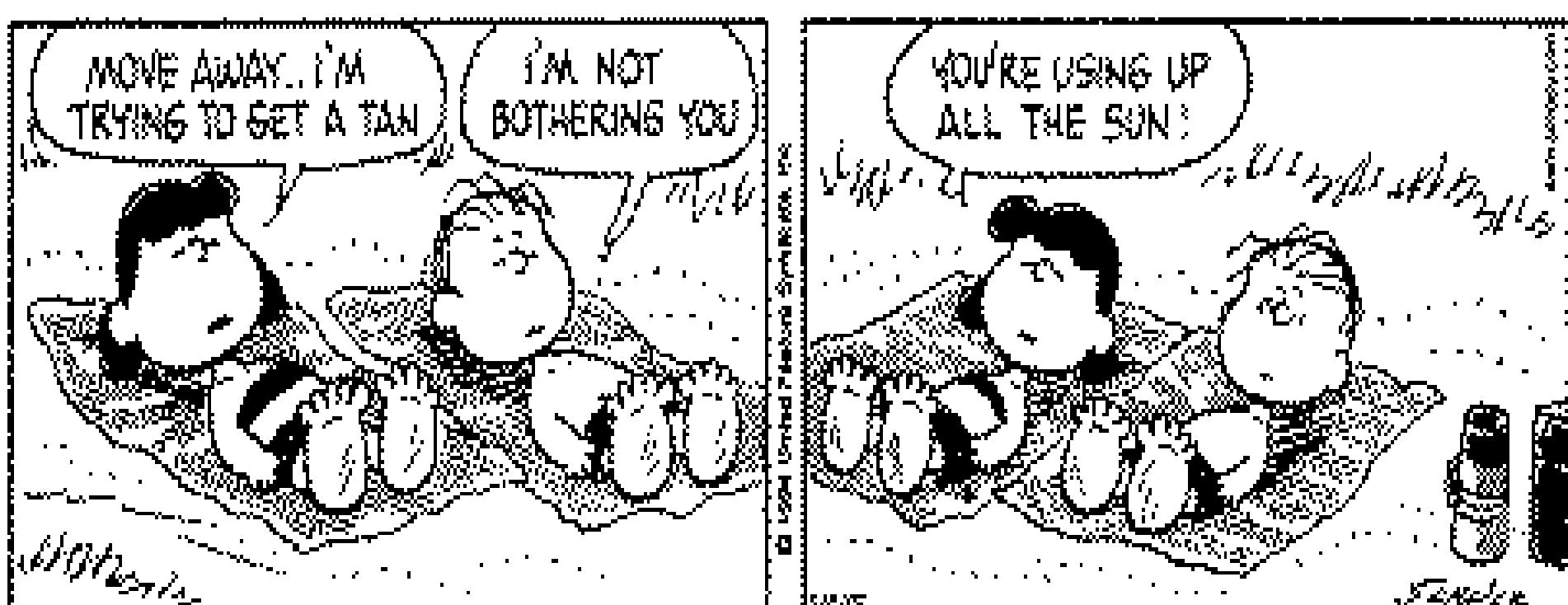
## Example: Cut and Choose



Player 1 cuts the cake in two equal pieces  
Player 2 chooses his favorite piece  
Player 1 takes the remainder

## Externalities

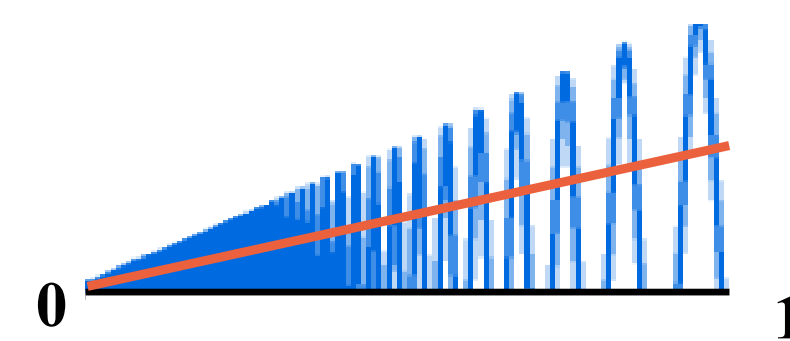
- Also known as **transaction spillovers**: third parties are influenced by transactions they did not agree to
  - > **Negative externalities**: pollution, smoking, risky choices (e.g. drinking and driving), overfishing
  - > **Positive externalities**: education, immunization, environmental cleanup, research (!)



- Common assumption in resource allocation: an agent's welfare is not affected by the consumption bundles of others (no externalities)
- *Envy is about externalities*: agents compare their own allocations with those of others.

## Externalities in Cake Cutting

- Generalized model to capture externalities when there are synergies between agents (e.g. altruism).
- Each agent has several value densities:  $V_{i,j}(X)$  is the value agent  $i$  gets from the allocation of piece  $X$  to agent  $j$
- The value of agent  $i$  under an allocation  $A$  is:  $V_i(A) = \sum_{j=1}^n V_{i,j}(A_j)$
- The optimal allocation of one agent can require infinitely many cuts:



## Fairness Criteria

### Proportionality:

- Allocation  $A$  is proportional if  $V_i(A_i) \geq 1/n$ ,  $\forall i \in N$

### Swap Envy-Freeness:

- Allocation  $A$  is swap envy-free if

$$V_{i,j}(A_i) + V_{i,j}(A_j) \geq V_{i,j}(A_j) + V_{i,j}(A_i), \forall i, j \in N$$

### Swap Stability:

- Allocation  $A$  is swap stable if

$$V_{i,j}(A_j) + V_{i,k}(A_k) \geq V_{i,j}(A_k) + V_{i,k}(A_j), \forall i, j, k \in N$$

## Properties

**Existence**: Swap-stability implies swap envy-freeness and proportionality; swap envy-freeness and proportionality are unrelated

**Lower bounds**: A swap envy-free and proportional allocation can require strictly more than  $n-1$  cuts

**Upper bounds**: Swap-stable allocations (which are also proportional and swap envy-free) are guaranteed to exist and require at most  $(n-1)n^2$  cuts when the value densities are continuous

## Query Model with Externalities

- Extended Robertson-Webb:
  - > **Cut** $_{i,j}(x, a)$ : Agent  $i$  returns  $y$  such that  $V_{i,j}([x, y]) = a$
  - > **Evaluate** $_{i,j}(x, y)$ : Return  $a$  such that  $V_{i,j}([x, y]) = a$
- An allocation that guarantees  $1/n^2$  to each agent can be computed with  $O(n^2)$  queries in the extended Robertson-Webb model
- **But there is no finite protocol that can compute a proportional allocation even for two agents in the extended Robertson-Webb model.**
- Alternative:
  - > **Cut Optimal** $(x, a)$ : Agent  $i$  outputs  $y$  such that  $i$ 's optimal allocation on  $[x, y]$ ,  $A_a$ , gives the agent exactly  $a$ :  $V_i(A_a) = a$
  - > **Evaluate Optimal** $_{i,j}(x, y)$ : Agent  $i$  outputs pair  $(A_a, a)$  such that  $A_a$  is an optimal allocation for  $i$  on  $[x, y]$  and  $V_i(A_a) = a$
- Every proportional protocol from the standard model translates to a proportional protocol with externalities when the Cut and Evaluate queries are replaced by Cut Optimal and Evaluate Optimal.